Modeling of TCP flows

Analysis of TCP

Image: A matrix

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# Dynamics and Stability of AIMD Algorithms

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The Hamilton Institute, NUI Maynooth

Hamilton Institute Workshop on Congestion Control September 27–28, 2005

based on joint work with: Abraham Berman, Christopher King, Douglas Leith, Robert Shorten, Rade Stanojević



# Modeling of TCP flows

The AIMD algorithm A linear model I A linear model II Model Validation

# Analysis of TCP

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# A Markov Model



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#### Many sources, one bottleneck



 $w_i(k) \stackrel{\circ}{=}$  the number of packages sent by the *i*-th source at the *k*-th congestion event.



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#### Time evolution for one source



 $w_i(k) \stackrel{\circ}{=}$  the number of packages sent by the *i*-th source at the *k*-th congestion event.



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#### Many sources, one bottleneck



 $w_i(k) \stackrel{\circ}{=}$  the number of packages sent by the *i*-th source at the *k*-th congestion event.



#### **Assumptions I**

- Congestion occurs in a single bottleneck.
- Congestion is noticed one RTT after it happens.
- Buffer size of bottleneck is small, i.e. RTT can be approximated by a constant.
- ▶ RTT is the same for all sources.
- The network is synchronized.



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#### Many sources, one bottleneck





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Analysis of TCP

#### A linear model

A little manipulation shows that

$$w(k+1) = Aw(k), \quad k \ge 1$$

where A is given by

$$\begin{bmatrix} \beta_1 & 0 & \dots & 0 \\ 0 & \beta_2 & & \vdots \\ \vdots & & \ddots & \\ 0 \dots & 0 & & \beta_n \end{bmatrix} + \frac{1}{\sum_i \alpha_i} \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \dots \\ \alpha_n \end{bmatrix} \begin{bmatrix} 1 - \beta_1 & \dots & 1 - \beta_n \end{bmatrix}.$$

Here  $\alpha_i$  is the additive increase parameter of the *i*-th source and  $\beta_i$  is the multiplicative decrease parameter.



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#### **Digression: Different RTT's**





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### **Digression: Different RTT's**

A little manipulation leads to a linear equation

$$w(k+1) = \widetilde{A}w(k)$$
.

The similarity transformation

$$\hat{w} := \begin{bmatrix} \operatorname{RTT}_{1}^{-1} & 0 & \dots & 0 \\ 0 & \operatorname{RTT}_{2}^{-1} & \dots & 0 \\ & & \ddots & \\ 0 & \dots & \operatorname{RTT}_{n}^{-1} \end{bmatrix} w$$

leads to an evolution equation for  $\hat{w}$  of the same form as in the case for equal RTT's.



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#### A linear model

$$w(k+1) = Aw(k), \quad k \ge 1$$

where *A* is given by

$$\begin{bmatrix} \beta_1 & 0 & \dots & 0 \\ 0 & \beta_2 & & \vdots \\ \vdots & & \ddots & \\ 0 \dots & 0 & & \beta_n \end{bmatrix} + \frac{1}{\sum_i \alpha_i} \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \cdots \\ \alpha_n \end{bmatrix} \begin{bmatrix} 1 - \beta_1 & \cdots & 1 - \beta_n \end{bmatrix}.$$

Here  $\alpha_i$  is the additive increase parameter of the *i*-th source and  $\beta_i$  is the multiplicative decrease parameter.



Analysis of TCP

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#### The Synchronized Case

Assumption: Every source experiences all congestions

$$A = \begin{bmatrix} \beta_1 & 0 & \dots & 0 \\ 0 & \beta_2 & & \vdots \\ \vdots & & \ddots & \\ 0 \dots & 0 & & \beta_n \end{bmatrix} + \frac{1}{\sum_i \alpha_i} \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \cdots \\ \alpha_n \end{bmatrix} \begin{bmatrix} 1 - \beta_1 & \cdots & 1 - \beta_n \end{bmatrix}.$$

The matrix A is positive and column stochastic. Thus by the Perron-Frobenius Theorem the evolution of  $A^k$  is very well understood.



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#### The Synchronized Case

$$A = \begin{bmatrix} \beta_1 & 0 & \dots & 0 \\ 0 & \beta_2 & & \vdots \\ \vdots & & \ddots & \\ 0 \dots & 0 & & \beta_n \end{bmatrix} + \frac{1}{\sum_i \alpha_i} \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \cdots \\ \alpha_n \end{bmatrix} \begin{bmatrix} 1 - \beta_1 & \cdots & 1 - \beta_n \end{bmatrix}.$$

**Theorem** (Berman, Shorten, Leith)

1. A has an eigenvalue one with eigenvector

$$\mathbf{v} = \begin{bmatrix} rac{lpha_1}{1-eta_1} & \cdots & rac{lpha_n}{1-eta_n} \end{bmatrix}$$

2. For all initial conditions  $w_0$  we have

$$A^k w_0 \to \theta v$$
.

3. the rate of convergence is exponential, bounds for this rate can be given in terms of the  $\beta_i$ .



#### **Assumptions II**

- Congestion occurs in a single bottleneck.
- Congestion is noticed one RTT after it happens.
- Buffer size of bottleneck is small, i.e. RTT can be approximated by a constant.
- RTT is the same for all sources.
- ► The network is synchronized.



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#### The unsynchronized case

$$w(k+1) = A(k)w(k), \quad k \geq 1$$

where A(k) is given by

$$\begin{bmatrix} \beta_1(k) & 0 & \dots & 0 \\ 0 & \beta_2(k) & \vdots \\ \vdots & & \ddots \\ 0 \dots & 0 & & \beta_n(k) \end{bmatrix} + \frac{1}{\sum_i \alpha_i} \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \cdots \\ \alpha_n \end{bmatrix} \begin{bmatrix} 1 - \beta_1(k) & \cdots & 1 - \beta_n(k) \end{bmatrix}$$

Here  $\alpha_i$  is the additive increase parameter of the *i*-th source and  $\beta_i(k)$  is equal to the multiplicative decrease parameter or 1, depending on whether the *i* - *th* source experiences congestion at the *k*-th congestion event or not.



#### The unsynchronized case

In the analysis of the dynamics of TCP flows we are led to the consideration of a linear inclusion of the form

$$w(k+1) \in \{Aw(k) \mid A \in \mathcal{M}\},\$$

where  $\mathcal{M} \subset \mathbb{R}^{n \times n}$  is the set of  $2^{n-1}$  matrices obtained by setting  $\beta_i, i = 1, \ldots, n$  either to 1 or to a constant in (0, 1). (Note that the identity is omitted.)



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#### **Model Validation I**



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#### Application to the TCP model

The matrices modeling TCP flows were of the form

$$\begin{bmatrix} \beta_1(k) & 0 & \dots & 0 \\ 0 & \beta_2(k) & \vdots \\ \vdots & & \ddots \\ 0 \dots & 0 & & \beta_n(k) \end{bmatrix} + \frac{1}{\sum_i \alpha_i} \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \cdots \\ \alpha_n \end{bmatrix} \begin{bmatrix} 1 - \beta_1(k) & \dots & 1 - \beta_n(k) \end{bmatrix}$$

The matrices are all column stochastic. The set  $\mathcal{M}$  is irreducible on the invariant subspace

$$S := \begin{bmatrix} 1 & 1 & \dots & 1 \end{bmatrix}^{\perp}$$



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#### Weak ergodicity

**Theorem** Let  $\{A(k)\}_{k\in\mathbb{N}} \subset \mathcal{M}^{\mathbb{N}}$  be a sequence with the property that at most one source does not see a drop infinitely often, then

$$\lim_{k\to\infty}A(k)A(k-1)\ldots A(0)_{|S}=0.$$

In other words, the sequence  $\{A(k)\}_{k\in\mathbb{N}}$  is weakly ergodic. This result extends results of Leizarowitz (LAA, 2000) in this special case. Proof builds on properties of extremal norms



#### A fairness result

Assume now that the matrices A(k) are i.i.d. random variables. These random variables induce for the *i*-th source a probability

 $\lambda_i := P($  source *i* experiences congestion at the *k*-th congestion event).

#### Theorem

If for each  $i = 1, \ldots, n$  we have

 $\lambda_i > 0$ ,

then, almost surely, (for the right  $\gamma$ )

$$\lim_{k\to\infty}\frac{1}{k+1}\sum_{l=0}^{k}w(l)=\gamma\begin{bmatrix}\frac{\alpha_1}{\lambda_1(1-\beta_1)}&\cdots&\frac{\alpha_n}{\lambda_n(1-\beta_n)}\end{bmatrix}^T$$



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#### Model Validation II, NS Simulation



Variation of mean of  $w_1(k)$ ,  $w_2(k)$  with propagation delay  $T_1$ . Key: +NS simulation result; · mathematical model; • Theorem on long time-averages; solid lines correspond to synchronised case.



#### Model Validation III, Data from a Real Network



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#### Model Validation III, Data from a Real Network

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Queue		rtt=21ms	rtt=34ms	rtt=55ms
20	Measurement	0.4055	0.3014	0.2931
	Theorem 3.3	0.4054	0.3061	0.2886
	% difference	0.0247	1.5594	1.5353
40	Measurement	0.4122	0.2849	0.3029
	Theorem 3.3	0.4121	0.2915	0.2964
	% difference	0.0243	2.3166	2.1459
60	Measurement	0.4024	0.3093	0.2882
	Theorem 3.3	0.4204	0.3087	0.2709
	% difference	4.4732	0.1940	6.0028
80	Measurement	0.416	0.3293	0.2547
	Theorem 3.3	0.3835	0.2891	0.3274
	% difference	7.8125	12.2077	28.5434



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Analysis of TCP

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#### A Markov model

We are interested in the system

$$W(k+1) = A(k)W(k), \quad A(k) \in \mathcal{M} := \{M_1, \ldots, M_{2^n-1}\}.$$

All matrices  $A_i$  are column stochastic, and window sizes are non-negative.

We may thus restrict our attention to the simplex

$$\Sigma := \left\{ x \in \mathbb{R}^n_+ \mid \sum_{i=1}^n x_i = 1 \right\} \,,$$



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We may thus restrict our attention to the simplex

$$\Sigma := \left\{ x \in \mathbb{R}^n_+ \mid \sum_{i=1}^n x_i = 1 \right\} \,,$$

Assume the random variables A(k), k = 0, 1, ... are i.i.d. and

$$P(A(k) = M_i) = \rho_i, \quad i = 1, \dots, \mu.$$

Assume also the "everyone drops" condition

$$\lambda_j = \sum_{\beta_j(M_i) < 1} \rho_i > 0 \,,$$



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The assumptions lead to a Markov process on  $\Sigma$ . The transition kernel of that process acts on continuous functions  $h: \Sigma \to \mathbb{R}$  through

$$Ph(x) = \int_{\Sigma} h(y)P(x, dy) = \sum_{i=1}^{\mu} \rho_i h(A_i x).$$

For each continuous h the sequence

$$P^kh$$
,  $k \in \mathbb{N}$ ,

is equicontinuous.

Markov chains with this property are called e-chains.



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**Theorem**(Meyn & Tweedie)

Consider an e-chain on a compact space  $\Sigma$ . If there is a positive and aperiodic state  $z \in \Sigma$ , then

- (i) there exists a unique invariant probability  $\pi$ ,
- (ii) for every  $x \in \Sigma$  and every continuous function  $h: \Sigma \to \mathbb{R}$  we have that if W(0) = x, then

$$\lim_{k\to\infty}\frac{1}{k}\sum_{j=0}^{k-1}h(W(j))=\int_{\Sigma}h(y)d\pi(y)\,,\quad\text{ in probability},$$

(iii) for every  $x \in \Sigma$  and every continuous function  $h : \Sigma \to \mathbb{R}$  we have

$$\int_{\Sigma} h(y) \mathcal{P}^k(x, dy) 
ightarrow \int_{\Sigma} h(y) d\pi(y) \,, \quad ext{ as } k 
ightarrow \infty \,.$$



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#### **Positive and Aperiodic States**

 $\mathcal{L} := \{ \{A_k\}_{k \in \mathbb{N}} \in \mathcal{M}^{\mathbb{N}} \mid \{A_k\}_{k \in \mathbb{N}} \text{ each source sees infinitely many drops} \}$ 

$$\mathcal{R}_{\mathcal{L}} := \{ R \mid \operatorname{rank} R = 1, \exists \{A_k\}_{k \in \mathbb{N}} \in \mathcal{L}, k_l \to \infty : \lim_{l \to \infty} \Pi(k_l) = R \}.$$

$$\mathcal{C} := \{ z \in \Sigma \mid zy^T \in \mathcal{R}_{\mathcal{L}} \}.$$



#### **Positive and Aperiodic States**

### Proposition

- (i)  $\ensuremath{\mathcal{C}}$  is compact and forward invariant,
- (ii) for any  $z \in C$  and any open neighborhood  $U \subset \Sigma$  of z there is a  $k_0 > 0$  such that  $P^k(x, U) > \delta > 0$ , for all  $k \ge k_0$  and all  $x \in \Sigma$ ,

(iii) For any initial condition  $\mathit{W}_0\in\Sigma$  we have almost surely

$$\lim_{k\to\infty}\operatorname{dist}(W(k),\mathcal{C})=0\,.$$

By (ii)& (iii): C is the set of positive, aperiodic states. C is the support of the unique invariant probability measure  $\pi$ .



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$$W(k+1) = A(k)W(k).$$

#### Theorem

Consider the Markov chain W on  $\Sigma$ . For this chain there is a positive and aperiodic state  $z \in \Sigma$ , and so

- (i) there exists a unique invariant probability  $\pi$ ,
- (ii) for every  $x \in \Sigma$  and every continuous function  $h : \Sigma \to \mathbb{R}$  we have that if W(0) = x, then

$$\lim_{k\to\infty}\frac{1}{k}\sum_{j=0}^{k-1}h(W(j))=\int_{\Sigma}h(y)d\pi(y)\,,\quad\text{ in probability},$$



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#### What does the support look like ?





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# Thank you !



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