

On buffering hypothesis in 802.11 analytic models

K. D. Huang and K. R. Duffy

Abstract—Using detailed statistical analysis of data taken from simulations and test-bed experiments, we have recently reported the inappropriateness of a queue-decoupling approximation that has been implicitly adopted by many distinct authors to incorporate queuing behavior in analytic models of finite load 802.11 networks. In the present paper we show that this flawed hypothesis leads to network throughput prediction errors in the presence of stations with asymmetric offered loads. We suggest that care should be taken in drawing deductions from models that adopt this hypothesis and that further modeling innovation is necessary.

Index Terms—WLAN, IEEE 802.11, Performance Evaluation.

I. INTRODUCTION

Bianchi’s seminal papers [1][2] introduced a celebrated mean-field Markov model of saturated WLANs employing IEEE 802.11. Assuming that each station in the network always has packets to send, this modeling approach enables the accurate prediction of the long run maximum stable throughput of an 802.11 infrastructure mode network. Bianchi’s methodology is based on a collision decoupling approximation. Assume that each station i has a fixed probability of collision given it is attempting transmission, p_i , that does not depend on collision history. Then the station’s back-off counter can be identified as an irreducible, aperiodic Markov chain from which the stationary probability that station i is attempting transmission, $\tau_i := \tau_i(p_i)$, can be determined explicitly as a function of p_i , as well as station and network parameters. The self-consistent fixed point equations

$$1 - p_i = \prod_{j \neq i} (1 - \tau_j(p_j)), \text{ for each station } i, \quad (1)$$

state that the probability station i experiences no collision given it is attempting transmission is the probability that no other station is attempting transmission. The unique solution $\{p_i^*\}$ of the system of equations (1) determines the operating point of the network. From $\{p_i^*\}$, fundamental performance quantities, such as per-station throughput, can be predicted.

Due to its intuitive appeal and its predictive accuracy, Bianchi’s approach has been widely adopted and developed for use in models that expand on its original range of applicability. For example, as stations in a typical network are not saturated, many authors have extended the paradigm to incorporate stations with finite load. For these models to be mathematically tractable, packet arrival at stations are assumed to follow a Poisson process. In these finite load models, if stations can buffer only a single packet, then no additional hypotheses beyond Bianchi’s original decoupling assumption

are necessary (e.g. [3][4][5]). These models continue to make accurate predictions.

To analytically model stations with packet buffers, two distinct approaches have been proposed: (A) the expansion of the original model’s state space to include buffer occupancy (e.g [6]); (B) the adoption of an additional queue-decoupling assumption, which we explain below. The drawback of the former approach is that the stationary distribution of the resulting Markov chain cannot be determined explicitly, so that identifying the functions $\tau_i(p_i)$ is numerically time consuming. Solving these models can take as long as the time to run a packet-level simulation. The potential difficulty with the latter methodology is that an additional approximation has been adopted and thus there is the possibility of new approximation errors. The advantage of that methodology is that $\tau_i(p_i)$ is, typically, determined as an explicit function, leading to simpler and faster numerics, giving these models a significant run-time advantage over simulations.

The queue-decoupling approximation, which is often implicitly assumed by authors, can be described as follows. Select a given station. After the k^{th} successful transmission, define $Q_k := 0$ if the station’s packet buffer is empty and $Q_k := 1$ if the station’s packet buffer is not empty. That is, $Q_k = 1$ if there is at least one packet awaiting transmission in the station’s buffer immediately after the k^{th} successful transmission and otherwise it is 0. The probability that $Q_k = 1$ is called the queue-busy probability. It is commonly assumed (e.g. [7][8][9][10]) that: (i) the sequence $\{Q_k\}$ consists of independent random variables; and (ii) the sequence $\{Q_k\}$ consists of identically distributed random variables that, in particular, do not depend on the back-off stage at which the most recent successful transmission was made. Rather than checking these assumptions directly, finite load models that adopt it are validated by comparing their predictions of throughput and delay with simulations, typically in symmetrically loaded networks. Using detailed statistical analysis of data from simulation [11] and experiment [12] we have recently shown, inter alia, that Bianchi’s collision-decoupling hypothesis is accurate even for non-saturated stations with buffers. However, while the queue-decoupling approximation makes mathematical analysis amenable, we have found that the identically distributed hypothesis, (ii) above, is inappropriate. Figure 7 in [11] illustrates the extreme dependence of queue-busy probability on back-off stage and the failure of the queue-decoupling hypothesis. This can be readily understood: the more collisions a packet experiences, the longer its MAC service time is and, therefore, it is more likely that there is another packet awaiting service at that station upon its successful transmission.

Our aim here is to show that even though models that adopt the queue-decoupling assumption have been reported to

Hamilton Institute, NUI Maynooth. SFI grant RFP-07-ENEF530. Aug. ’09: corrected typos in eq. (2) and $E(B)$. These are consistent with the expressions used to create the figures in the published paper.

make accurate predictions, caution is needed in using their predictions. We do this by considering a setting where a queuing model can be built with or without the queue-decoupling approximation. In particular, we consider a model where stations can buffer one packet beyond the MAC. This enables us to show that even though the models make remarkably similar throughput predictions for symmetrically loaded networks, that under asymmetric loads, which are common in practice, they disagree.

II. MODELS

The basic structure of the models introduced here is similar to nearly all of those developed from Bianchi's paradigm. Model specifics were chosen as they form the simplest situation where an analytic model can be developed with or without the queue-decoupling hypothesis. Bianchi's approach is well established, so we present only key equations rather than their derivation. As in most analytic work, we assume clear channel conditions so that transmission failure is only caused by collisions and that there are no hidden stations. We also assume that there is no post back-off, so that every packet experiences at least one back-off period.

Consider a network of stations labeled $\{1, \dots, N\}$. Let station i have Poisson arrival rate λ_i , minimum contention window W_i , maximum contention window $2^{m_i} W_i$ and assume that packets are discarded after M_i ($\geq m_i$) collisions. Adopting Bianchi's fundamental hypothesis, the conditional collision probability given attempted transmission p_i is assumed not to depend on back-off stage. In addition, assume that each station can buffer a single packet in addition to the one being processed by the MAC. Temporarily suppressing the station subscript i , each station's back-off process is described by the Markov chain depicted in Figure 1, where the state (k, l) corresponds to the current packet in the MAC having experienced k collisions and having a back-off counter with value l . The Idle state corresponds to the station having no packet awaiting transmission. The Markov chain has the following parameters: r , the probability that at least one packet arrives while the station is in the idle state, p the conditional collision probability and $\{q_j\}$ the queue-busy probability at back-off stage j . We will relate the r and $\{q_j\}$ parameters to the Poisson packet arrival rate λ , depending on whether the queue decoupling assumption is adopted or not. In the general setting the stationary distribution, $b(l, k)$, of this Markov chain can be determined explicitly as a function of these parameters. In particular, defining the normalization

$$\frac{1}{\eta} := \frac{1}{2} \left(\frac{1-p-p(2p)^m}{(1-p)(1-2p)} W - \frac{2^m p^{M+1}}{1-p} W - \frac{1-p^{M+1}}{1-p} \right) + \frac{1-p}{r} \left(\sum_{k=0}^{M-1} p^k (1-q_k) + \frac{p^M}{1-p} (1-q_M) \right), \quad (2)$$

the stationary probability of attempted transmission is

$$\tau(p, r, \{q_j\}) = \frac{1-p^{M+1}}{1-p} \eta. \quad (3)$$

In order to relate the offered load parameters r and $\{q_j\}$ we consider two models, one with and one without the queue-decoupling approximation.

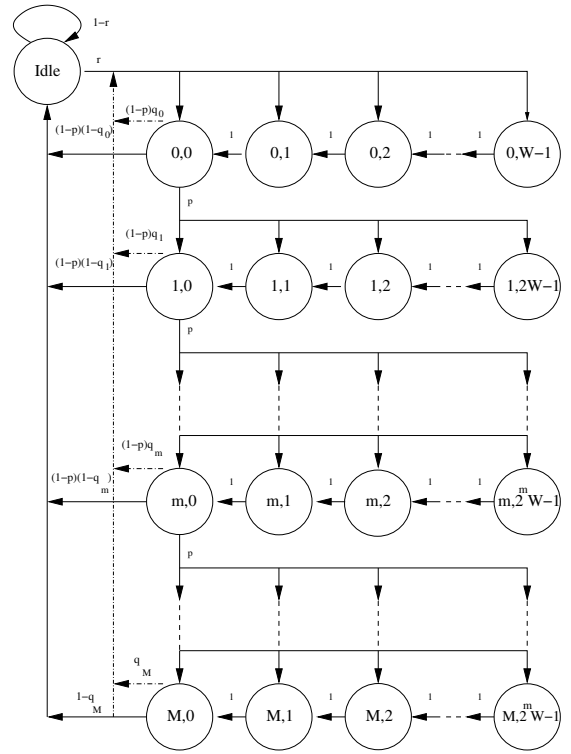


Fig. 1. Station back-off counter Markov chain

Model (I): Const- q . Adopting the queue-decoupling approximation, $q_j := q$ for all j and $1-q$ is the probability that no packet arrives during an average MAC service time. With idle slots being of length δ ($20\mu s$ in 802.11b) and busy slots of length T_b (for ease of presentation, we assume that collisions and successful transmissions take the same time; $578\mu s$ plus payload at 11Mbps in 802.11b with ACK header at 1Mbps), the expected time between counter decrements is

$$T := \prod_{i=1}^N (1-\tau_i) \delta + \left(1 - \prod_{i=1}^N (1-\tau_i) \right) T_b,$$

while the average time from when a packet starts being processed by the MAC until it successful transmission is

$$E(B) = \frac{W(1-p-p(2p)^m)}{2(1-2p)(1-p)} + \frac{p^{M+1}(1-W2^m)}{2(1-p)}.$$

Based on the reasoning in [5], we use the standard approximations that $1-r = \exp(-\lambda T)$ and $1-q = \exp(-\lambda E(B)T) = (1-r)^{E(B)}$, so that the final term in equation (2) reduces to $(1-r)^{E(B)}/r$. To solve the Const- q model, using (3), one searches for a solution of fixed point equations (1).

Model (II): Var- q . For the model without the queue-decoupling approximation, we wish to assume as little as possible beyond Bianchi's collision decoupling hypothesis. This leads us to treating slots being idle or busy as an i.i.d. process as well as having the queue-busy probability depend on backoff stage. Again suppressing station-based subscripts,

we have that

$$1 - r = \prod_{k \neq i} (1 - \tau_k) e^{-\lambda \delta} + \left(1 - \prod_{k \neq i} (1 - \tau_k) \right) e^{-\lambda T_b}.$$

For $0 \leq j < m$, we have that

$$1 - q_j = \left(\prod_{l=0}^j \frac{1}{2^l W} \right) e^{-\lambda T_b(j+1)} \left(\prod_{l=0}^j \frac{1 - (1-r)^{2^l W}}{r} \right)$$

and, for $m \leq i \leq M$, we have the equality

$$1 - q_j = (1 - q_{j-1}) \frac{1 - (1-r)^{2^m W}}{2^m W r} e^{-\lambda T_b}.$$

To solve the Var-q model, using (3), one searches for a solution of these equations in addition to the fixed point equation (1).

III. RESULTS

Consider an 802.11b infrastructure mode network with N stations that have 802.11b parameters identical to those in Table I [5]. All packets have a 1000 byte payload. A model validation scenario of the sort commonly used by authors is where there are 10 stations and the arrival rates of all stations are the same. In Figure 2, model predictions of station throughputs are compared with the output of ns-2 packet level simulations. It can be seen that the Const-q and Var-q models are nearly inseparable and both make remarkably accurate throughput predictions when compared with results from simulations. Validation scenarios of this sort have led many authors to deduce that the queue-decoupling hypothesis is a good approximation. However, Figure 3 reports on the same network, but where the arrival rate at one station is thirty times that of the nine others: $\lambda_1 = 30\lambda_2$. The predictions of the two models no longer coincide, with the Var-q model continuing to make accurate predictions, while the Const-q model does not. Other asymmetric slices through the throughput versus offered load surface for networks of various sizes, N , show similar differences in predictions. This indicates that the flawed queue-decoupling hypotheses does follow through to prediction inaccuracies.

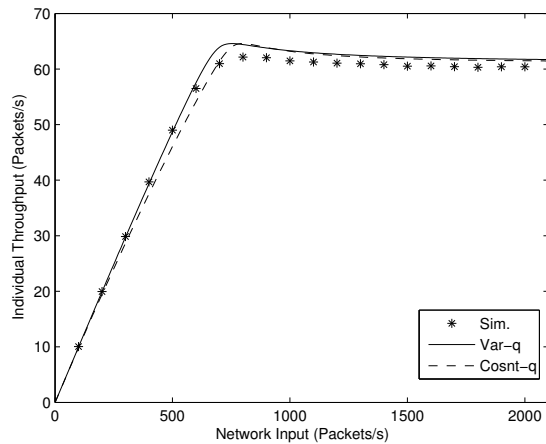


Fig. 2. Symmetric offered load. $N = 10$. Predictions and ns-2 simulation.

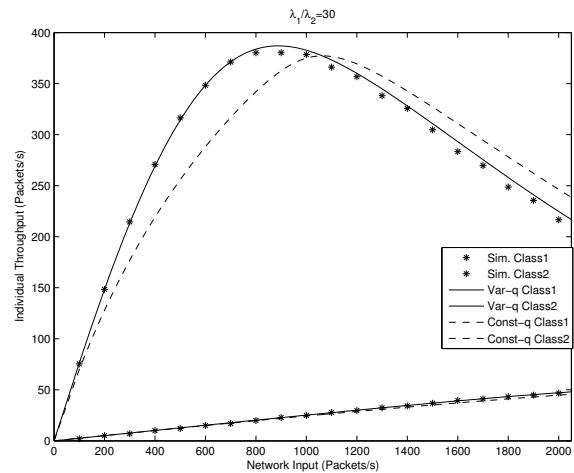


Fig. 3. Asymmetric offered load. $N = 10$. One station with arrival rate λ_1 and nine with rate λ_2 . Predictions and ns-2 simulation.

IV. CONCLUSIONS

We have shown how a commonly adopted, but flawed, hypothesis used in modeling 802.11 infrastructure model networks can lead to predictive errors. As finite load models are used, for example, to propose load-dependent 802.11e parameterizations to achieve fair bandwidth allocation in infrastructure and mesh networks (e.g. [13], [14]), these inaccurate predictions could lead to settings that perform poorly. This work suggests that to build an analytic model that incorporates buffering, is tractable and has run-time advantages over simulation, one must either accept the inaccuracies that come with the queue-decoupling hypothesis or the research community must rise to the challenge of providing new modeling innovation.

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